The chromatic number of a random lift of a regular graph

Xavier Pérez-Giménez joint work with JD Nir

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Probabilistic Combinatorics Online, Sep 2020

Lifts

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Surjective graph homomorphism $\Pi: L \to G$ that is a bijection between edges incident with ν and edges incident with $\Pi(\nu)$.

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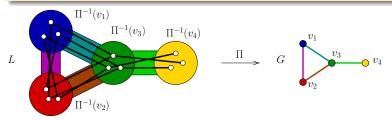
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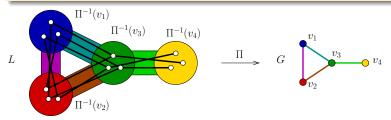


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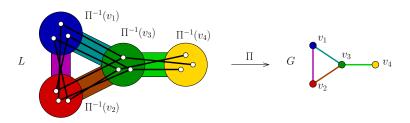
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Note: G may have loops and multiple edges.

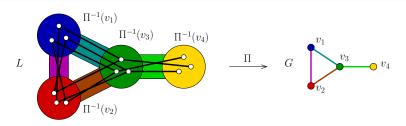
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Fact 1: G connected \Longrightarrow all **fibers** $\Pi^{-1}(v)$ have same cardinality.

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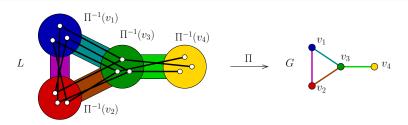


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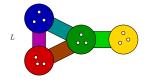
Fact 3: $\chi(L) \leq \chi(G)$





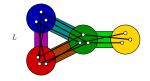
Random *n*-lift model (Amit, Linial 2002):

- Replace $v \in V(G)$ by bin with n vertices.
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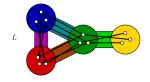
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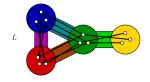
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Facts (for d-regular G):

- $G = B_{d/2} \Longrightarrow L$ contiguous to uniform d-regular multigraph.
- Not true for $G = K_{d+1}$



(Amit, Linial, Matoušek 2002)

Problem 1

For any
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: is $\chi(L) = \Omega\left(\frac{\chi(G)}{\log \chi(G)}\right)$ a.a.s.?

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Conjecture

For any G, there is k_G with $\chi(L) = k_G$ a.a.s.

...and many more open questions!



Our results

Thm (P-G, Nir 2019++):

Let $d \ge 3$ and $k_d = \min\{k \in \mathbb{N} : d < 2k \log k\}$. $(k_d \approx \frac{d}{2 \log d})$

Let L be random n-lift of $G = K_{d+1}$.

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Thm (Kemkes, P-G, Wormald 2010):

Analogous result holds for uniform d-regular graphs.

(Improved by Coja-Oghlan, Efthymiou, Hetterich 2016.)

Main tools

- Small subgraph conditioning method (Robinson & Wormald 1992)
- Optimization over stochastic matrices (Achlioptas, Naor 2005)
- Laplace summation method (Greenhill, Janson, Ruciński 2010)
- Saddle-point method
- Algebraic graph theory
 - Kirchhoff Matrix-Tree Thm
 - Counting non-backtracking closed walks

Proof structure

Lower bound on $\chi(L)$:

X = # k-colourings of L.

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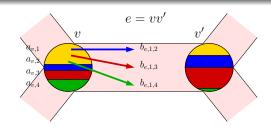
Thm: If $k > k_d$, then $EY^2 = \Theta((EY)^2)$

Then $P(Y > 0) \ge \frac{(EY)^2}{EY^2} \sim C > 0$ (Paley-Zygmund)

Unfortunately, C < 1 due to the influence of short cycles in L.

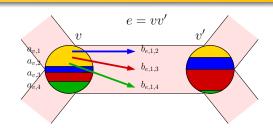


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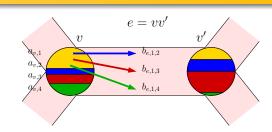
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Note: $(a_{v,i})$ is stochastic $|V| \times k$ matrix.

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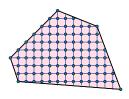
Note: $(a_{V,i})$ is stochastic $|V| \times k$ matrix.

Claim:

Max contribution is from $a_{v,i} = 1/k$, $b_{e,i,i'} = 1/k(k-1)$.

(We extend result by Achlioptas, Naor 2005.) 990

Summation domain:



$$\begin{cases} \mathbf{x} \in \mathcal{P} \\ B\mathbf{x} = \mathbf{y} \\ \mathbf{x} \in \left(\frac{1}{n}\mathbb{Z}\right)^D \end{cases}$$

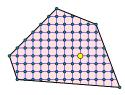
where:

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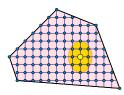
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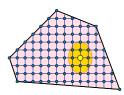
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(Laplace summ. / Saddlepoint method)

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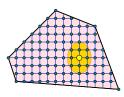
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• About C: It depends on

 $\begin{cases} \text{Hessian of } f \\ \text{Volume of fundamental cell in lattice} \end{cases}$

(Greenhill, Janson, Ruciński 2010)

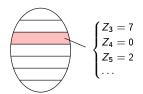
Instead, we count maximal forests in Γ with incidence matrix B.

Small subgraph conditioning (Robinson, Wormald 1992)

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 $Z_i = \#$ cycles of length i in L.

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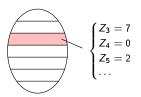


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Rough idea:

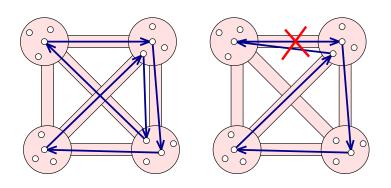
Suppose:

- $ullet rac{\mathsf{E}(\mathit{YZ}_i)}{\mathsf{E}\mathit{Y}} = 1 + \delta_i + o(1) \quad (\& \ \mathsf{joint factorial moments})$ (i.e. $\mathit{Z}_i \sim \mathsf{Poi}(1 + \delta_i)$ in space "weighted" by Y).
- $\bullet \ \frac{\mathsf{E}Y^2}{(\mathsf{E}Y)^2} = \exp(\sum_i \lambda_i \delta_i^2) + o(1).$

Then $P(Y > 0) \rightarrow 1$ (+ contiguity [...]).

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Counting non-backtracking closed walks



Some algebraic tools: (Friedman 2008)

(Amit, Linial, Matoušek 2002)

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Conjecture

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